

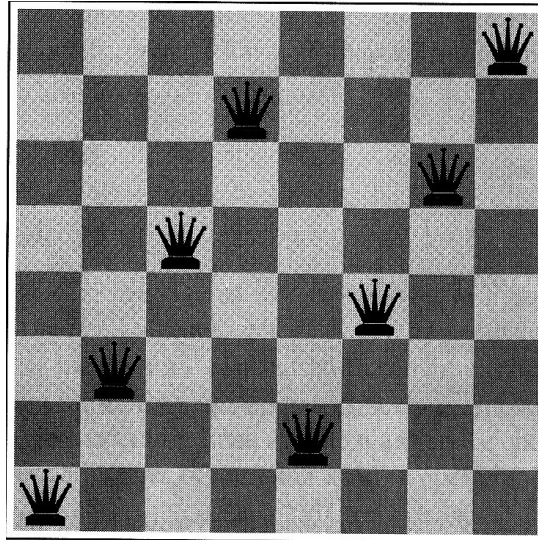


# Artificial Intelligence

## **Lecture:**

Problem Solving using Search - (Single agent search)  
Uninformed Search

# Example: 8-queens



- State: any arrangement of up to 8 queens on the board
- Operation: add a queen (incremental), move a queen (fix-it)
- Initial state: no queens on board
- Goal state: 8 queens, with no queen is attacked
- Solution Path: The set of operations that allowed you to get to the board that you see above at the indicated positions.

## Example: 8-puzzle

5	4	
6	1	8
7	3	2

Start State

1	2	3
8		4
7	6	5

Goal State

- State:
- Operators:
- Goal test:
- Solution path:

# Example: 8-puzzle

5	4	
6	1	8
7	3	2

Start State

1	2	3
8		4
7	6	5


Goal State

- **Operators:** moving blank left, right, up, down (ignore jamming)
- **Goal:** goal state
- **State:** integer location of tiles (ignore intermediate locations)
- **Solution:** move 4 tile to blank, move 1 tile blank, etc.



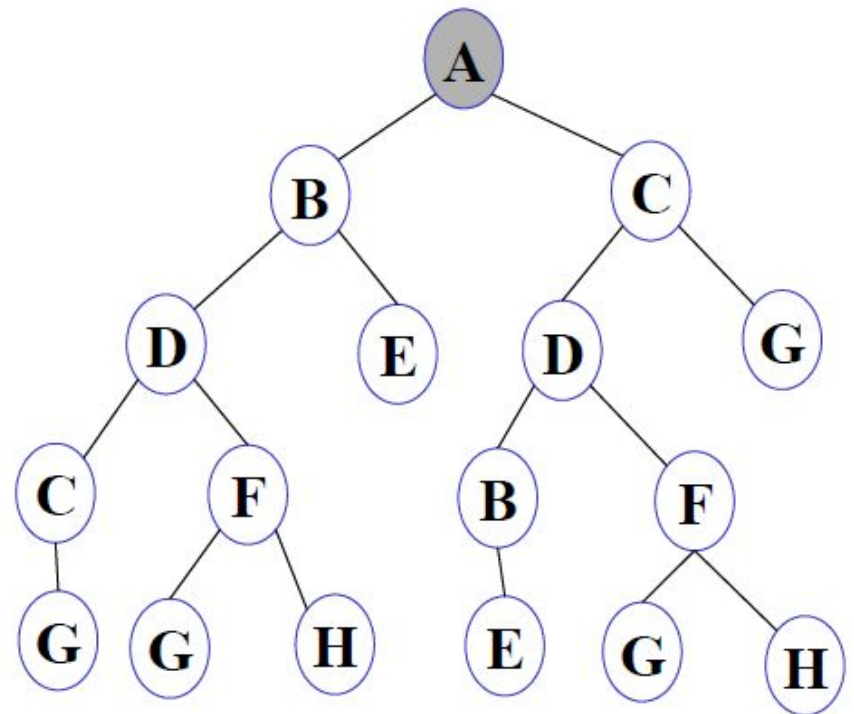
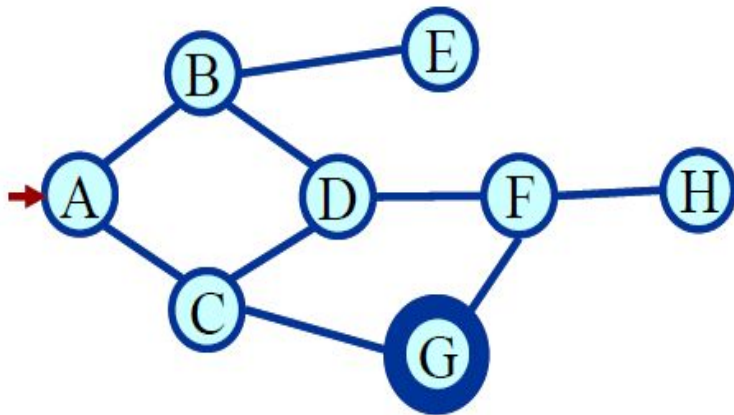
# Search

- What is Search?

- 
- Formulate appropriate problems in optimization and planning (sequence of actions to achieve a goal) as search tasks: initial state, operators, goal test, path cost

# Search Tree

List all Possible Path  
Eliminate Cycles from paths  
Result: Search Tree



# The basic search algorithm

```
Let L be a list containing the initial state (L= the fringe)
Loop
  if L is empty return failure
  Node ← select (L)
  if Node is a goal
    then return Node
      (the path from initial state to Node)
  else generate all successors of Node, and
    merge the newly generated states into L
End Loop
```



# Evaluating Search Strategies

- **Completeness**

- Guarantees finding a solution whenever one exists

- **Time complexity**

- How long (worst or average case) does it take to find a solution? Usually measured in terms of the number of nodes expanded

- **Space complexity**

- How much space is used by the algorithm? Usually measured in terms of the maximum size of the “nodes” list during the search

- **Optimality/Admissibility**

- If a solution is found, is it guaranteed to be an optimal one? That is, is it the one with minimum cost?



# Problem solving

- We want:
  - To automatically solve a problem
  
- We need:
  - A representation of the problem
  - Algorithms that use some strategy to solve the problem defined in that representation

# Problem representation

- General:
  - **State space:** a problem is divided into a set of resolution steps from the initial state to the goal state
  - **Reduction to sub-problems:** a problem is arranged into a hierarchy of sub-problems
- Specific:
  - Game resolution
  - Constraints satisfaction

# States

- A problem is defined by its elements and their relations.
- In each instant of the resolution of a problem, those elements have specific descriptors (How to select them?) and relations.
- A **state** is a representation of those elements in a given moment.
- Two special states are defined:
  - **Initial state** (starting point)
  - **Final state** (goal state)

## State modification: successor function

- A successor function is needed to move between different states.
- A **successor function** is a description of possible actions, a set of operators. It is a transformation function on a state representation, which convert it into another state.
- The successor function defines a relation of accessibility among states.
- Representation of the successor function:
  - Conditions of applicability
  - Transformation function



## State space

- The **state space** is the set of all states reachable from the initial state.
- It forms a graph (or map) in which the nodes are states and the arcs between nodes are actions.
- A **path** in the state space is a sequence of states connected by a sequence of actions.
- The solution of the problem is part of the map formed by the state space.

## Problem solution

- A **solution** in the state space is a path from the initial state to a goal state or, sometimes, just a goal state.
- **Path/solution cost**: function that assigns a numeric cost to each path, the cost of applying the operators to the states
- Solution quality is measured by the path cost function, and an **optimal solution** has the lowest path cost among all solutions.
- Solutions: any, an **optimal one**, all. Cost is important depending on the problem and the type of solution sought.

# Problem description

- Components:
  - State space (explicitly or implicitly defined)
  - Initial state
  - Goal state (or the conditions it has to fulfill)
  - Available actions (operators to change state)
  - Restrictions (e.g., cost)
  - Elements of the domain which are relevant to the problem (e.g., incomplete knowledge of the starting point)
  - Type of solution:
    - Sequence of operators or goal state
    - Any, an optimal one (cost definition needed), all



# Uninformed vs. informed search

- **Uninformed search strategies**

- Also known as “blind search,” uninformed search strategies use **no information about the likely “direction” of the goal node(s)**
- Uninformed search methods: Breadth-first, depth-first, depth-limited, uniform-cost, depth-first iterative deepening, bidirectional

- **Informed search strategies**

- Also known as “**heuristic search**” informed search strategies **use information about the domain to** (try to) (usually) head in the general direction of the goal node(s)
- Informed search methods: Hill climbing, best-first, greedy search, beam search, A, A\*



## Uninformed

1. Search without information
2. No Knowledge
3. Time Consuming
4. More complexity - time and space
5. BFS, DFS

## Informed

1. Search with information
2. use knowledge to find steps to solution
3. quick solution
4. less complexity
5. A\*, Best First, AO\*

# Breadth-First Search (BFS)

- Uninformed Search technique
- FIFO (Queue)
- Expand shallowest unexpanded node
- Level search
- Fringe: nodes waiting in a queue to be explored, also called **OPEN**
- **Optimal - shortest path**
- Implementation:
  - For BFS, *fringe* is a first-in-first-out (FIFO) queue
  - new successors go at end of the queue
- Repeated states?
  - Simple strategy: do not add parent of a node as a leaf

# BFS Algorithm

## **Breadth first search**

Let *fringe* be a list containing the initial state

Loop

    if *fringe* is empty return failure

    Node  $\leftarrow$  remove-first (*fringe*)

    if Node is a goal

        then return the path from initial state to Node

    else generate all successors of Node, and

        (merge the newly generated nodes into *fringe*)

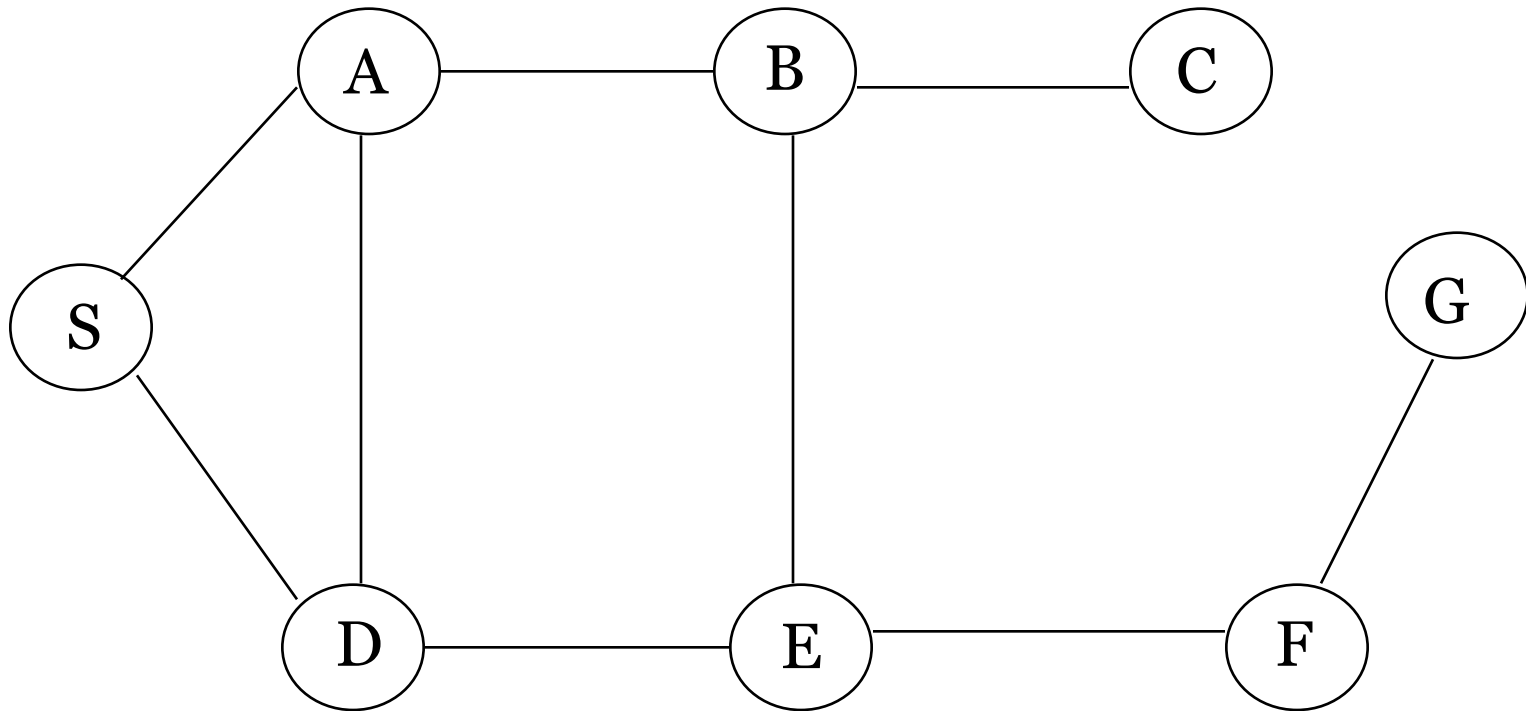
        add generated nodes to the back of *fringe*

End Loop

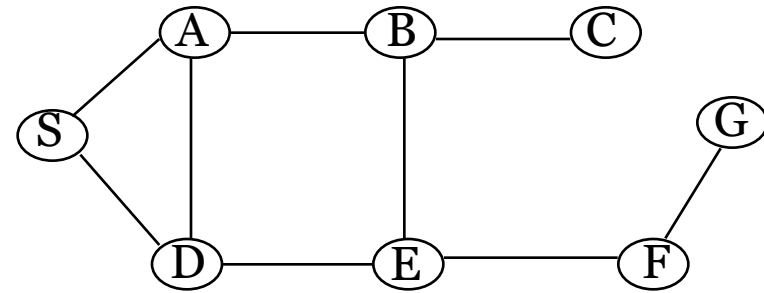
# Example: Map Navigation

State Space:

S = start, G = goal, other nodes = intermediate states, links = legal transitions



# BFS Search Tree



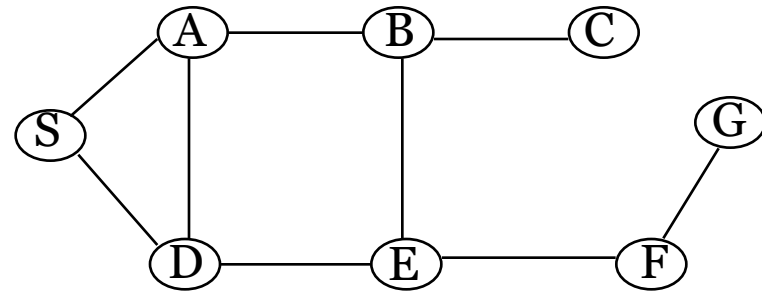
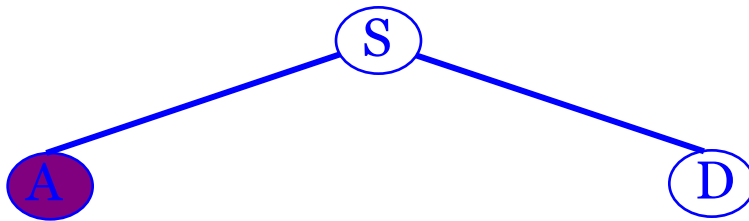
Queue = {S}

Select S

Goal(S) = true?

If not, Expand(S)

# BFS Search Tree



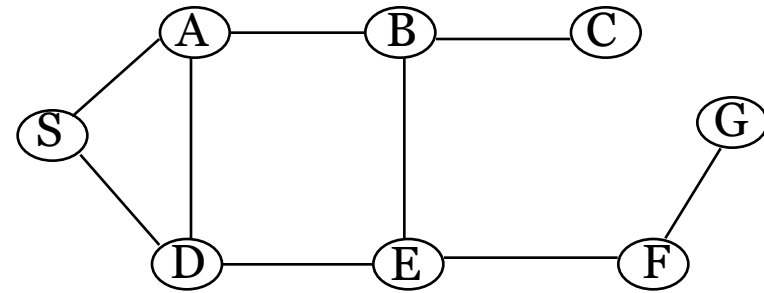
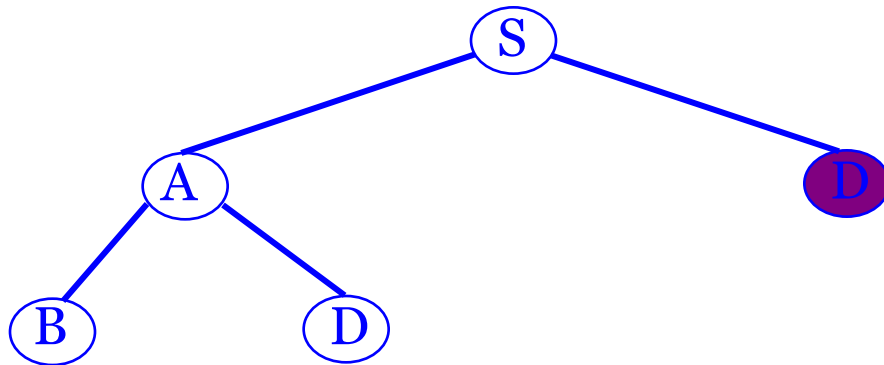
Queue = {A, D}

Select A

Goal(A) = true?

If not, Expand(A)

# BFS Search Tree



Queue = {D, B, D}

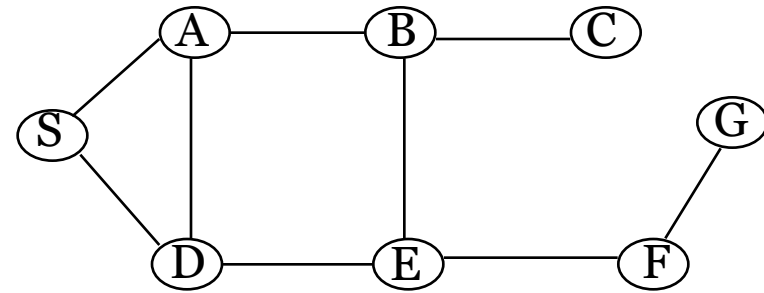
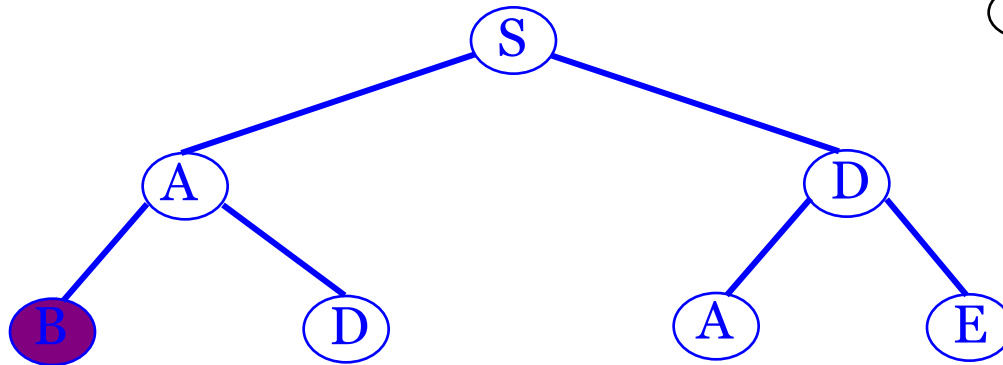
Select D

Goal(D) = true?

If not, expand(D)



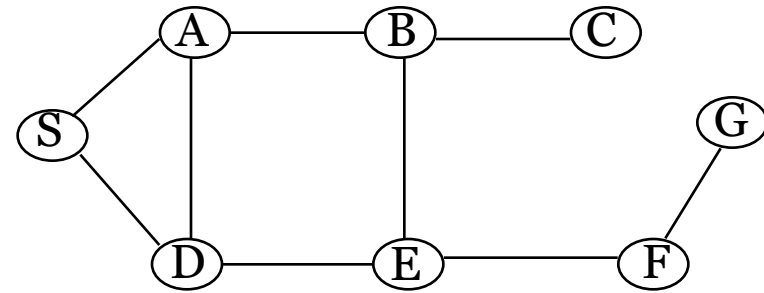
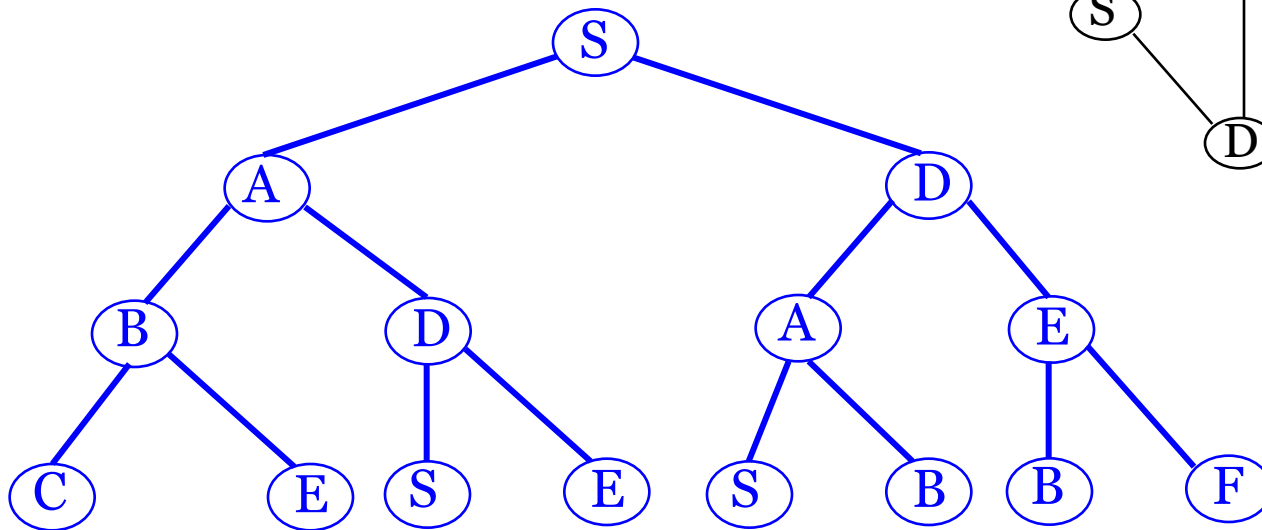
# BFS Search Tree



Queue = {B, D, A, E}

Select B  
etc.

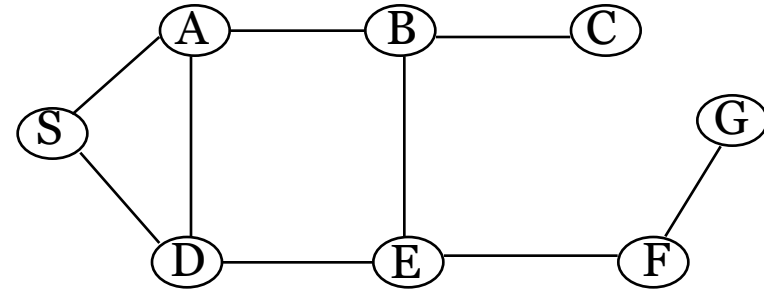
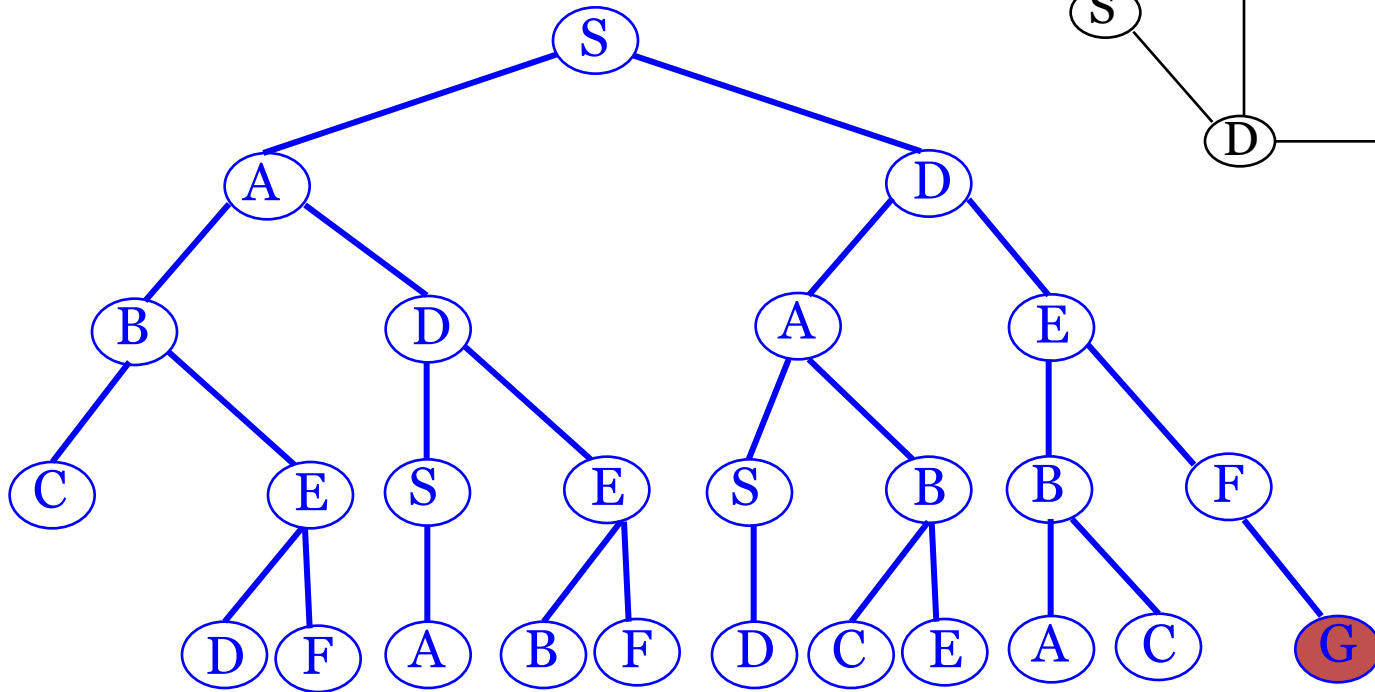
# BFS Search Tree



Level 3

Queue = {C, E, S, E, S, B, B, F}

# BFS Search Tree



Level 4  
Expand queue until G is at front  
Select G  
Goal(G) = true



# Depth-First Search (DFS)

- Uninformed Search
- Stack (LIFO)
- Expand deepest unexpanded node
- LIFO
- Deepest Node
- Incomplete ---> 1. Loop 2. search space infinite i.e. unlimited depth
- Implementation:
  - For DFS, *fringe* is a LIFO queue
  - new successors go at beginning of the queue
- Repeated nodes?
  - Simple strategy: Do not add a state as a leaf if that state is on the path from the root to the current node

Non Optimal  
Time complexity

# DFS Algorithm

## Depth First Search

Let *fringe* be a list containing the initial state

Loop

if *fringe* is empty return failure

Node ← remove-first (*fringe*)

if Node is a goal

then return the path from initial state to Node

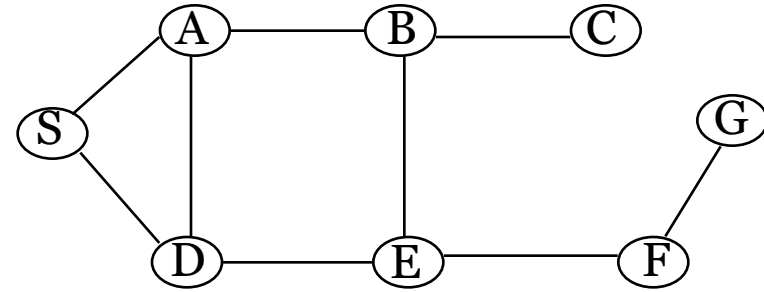
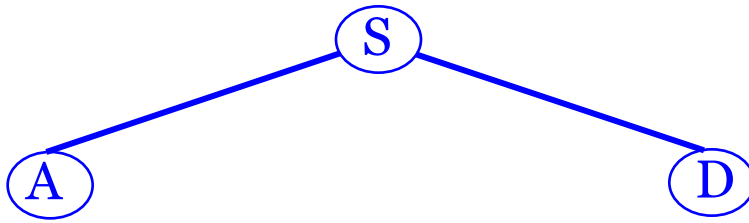
else generate all successors of Node, and

merge the newly generated nodes into *fringe*

add generated nodes to the front of *fringe*

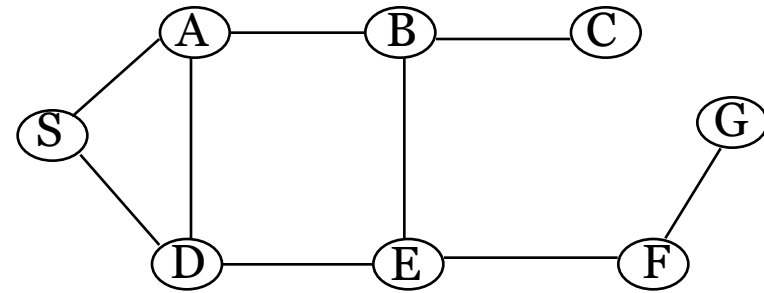
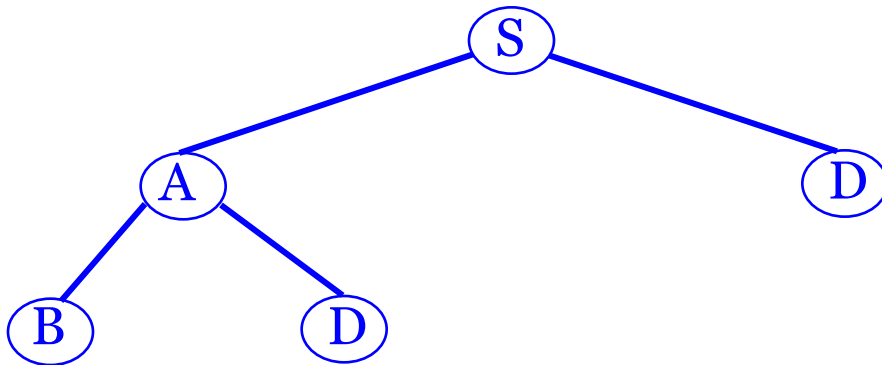
End Loop

# DFS Search Tree



Queue = {A,D}

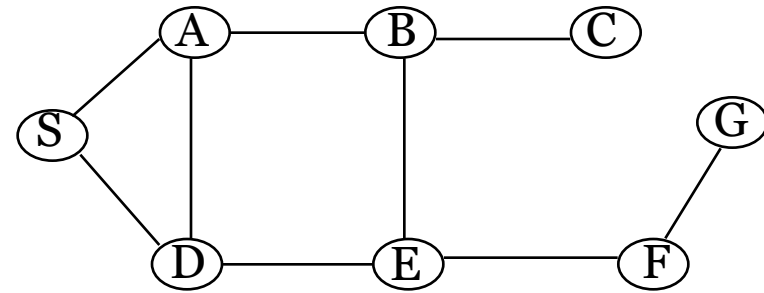
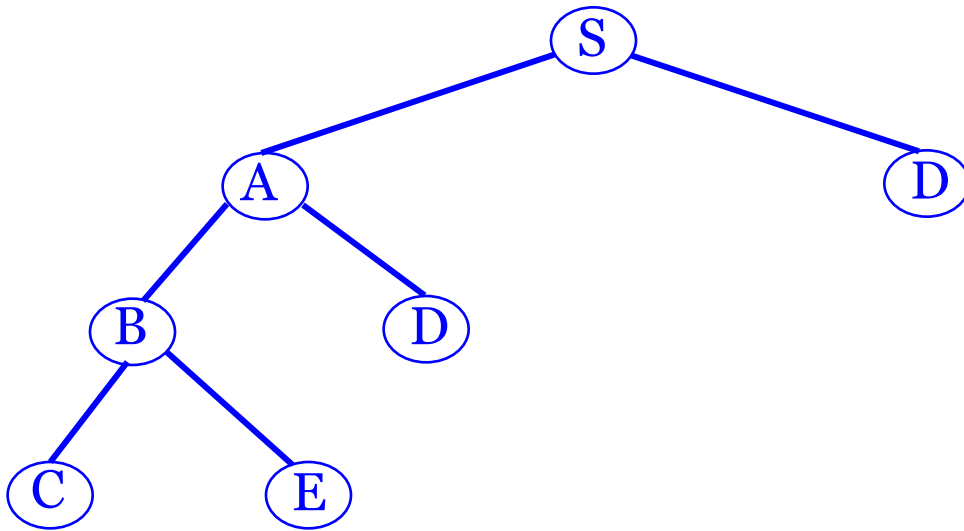
# DFS Search Tree



Queue = {B,D,D}

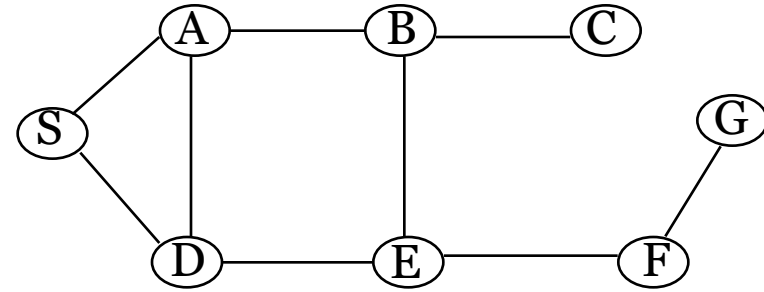
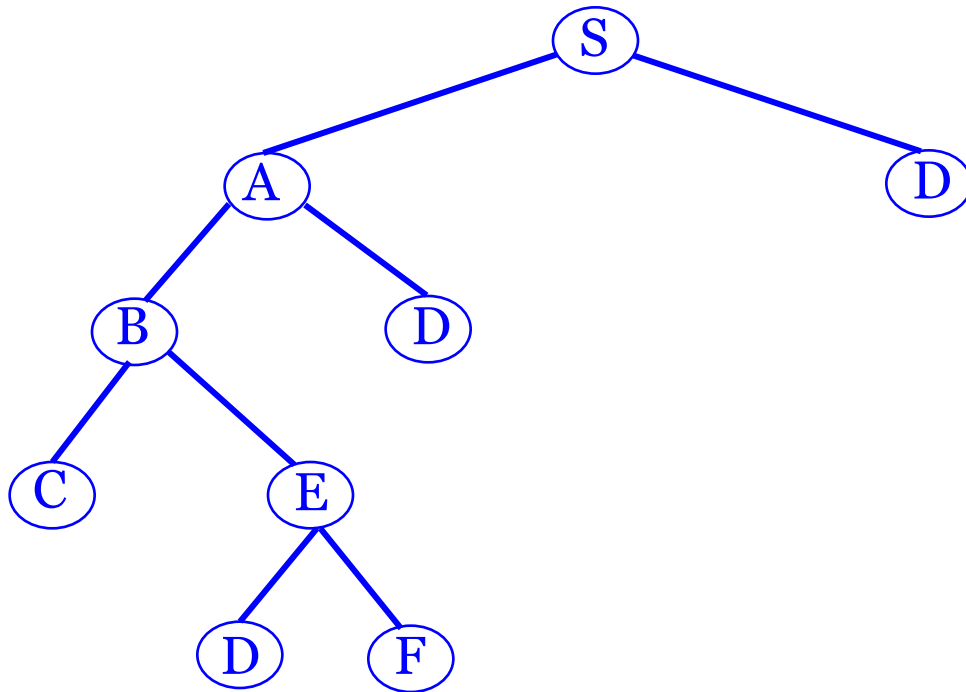


# DFS Search Tree



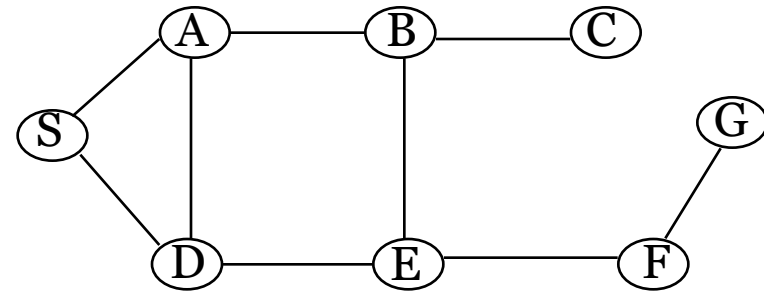
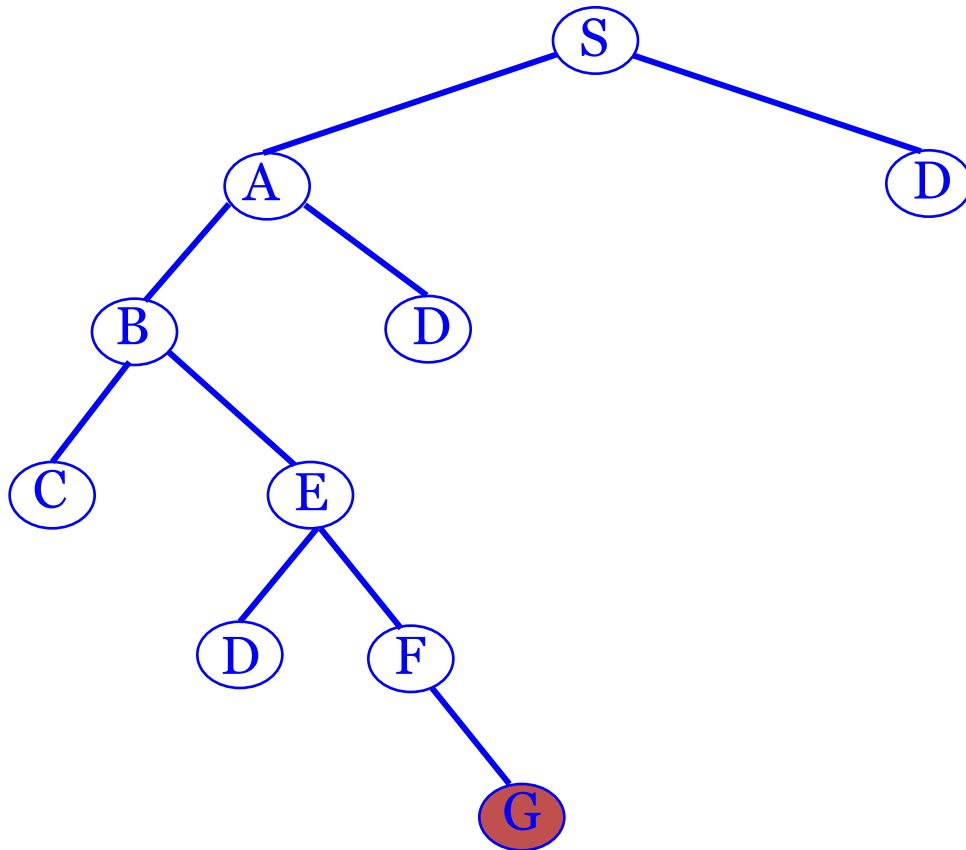
Queue = {C,E,D,D}

# DFS Search Tree



Queue = {D,F,D,D}

# DFS Search Tree



Queue = {G,D,D}



# Evaluation of Search Algorithms

- **Completeness**
  - does it always find a solution if one exists?
- **Optimality**
  - does it always find a least-cost (or min depth) solution?
- **Time complexity**
  - number of nodes generated (worst case)
- **Space complexity**
  - number of nodes in memory (worst case)
- **Time and space complexity are measured in terms of**
  - $b$ : maximum branching factor of the search tree
  - $d$ : depth of the least-cost solution
  - $m$ : maximum depth of the state space (may be  $\infty$ )

# Breadth-First Search (BFS) Properties

- Complete? **Yes**
- Optimal? **Yes**, for the shortest path
- Time complexity?  **$O(b^d)$**

$$1 + b + b^2 + \dots + b^d = O(b^d)$$

**exponential in the depth of the solution  $d$**

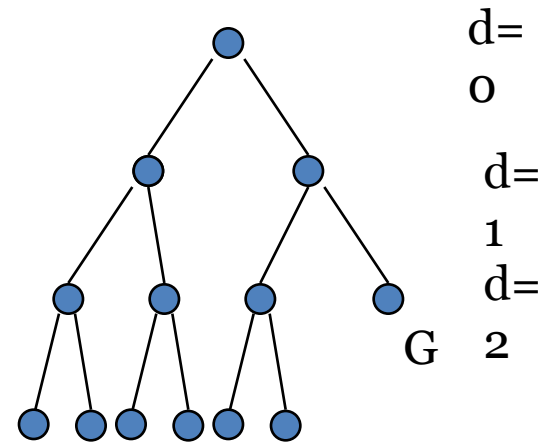
- Space complexity?  **$O(b^d)$**   
same as time - every node is kept in the memory
- **Main practical drawback?** exponential space complexity

# Complexity of Breadth-First Search

- **Time Complexity**

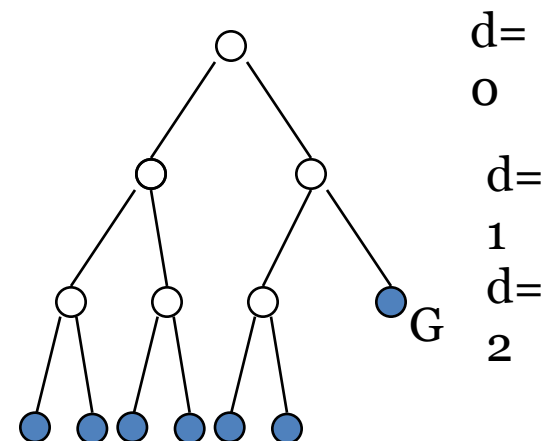
- assume (worst case) that there is 1 goal leaf at the RHS at depth  $d$
- so BFS will generate

$$= b + b^2 + \dots + b^d + b^{d+1} - b$$
$$= O(b^{d+1})$$



- **Space Complexity**

- how many nodes can be in the queue (worst-case)?
- at depth  $d$  there are  $b^{d+1}$  unexpanded nodes in the Q =  $O(b^{d+1})$



## Examples of Time and Memory Requirements for Breadth-First Search

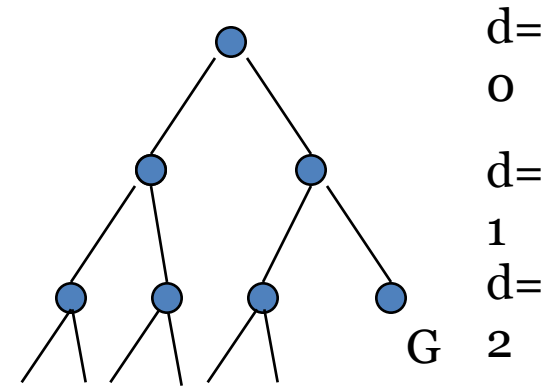
Assuming  $b=10$ , 10000 nodes/sec,  
1kbyte/node

Depth of Solution	Nodes Generated	Time	Memory
2	1100	0.11 seconds	1 MB
4	111,100	11 seconds	106 MB
8	$10^9$	31 hours	1 TB
12	$10^{13}$	35 years	10 PB

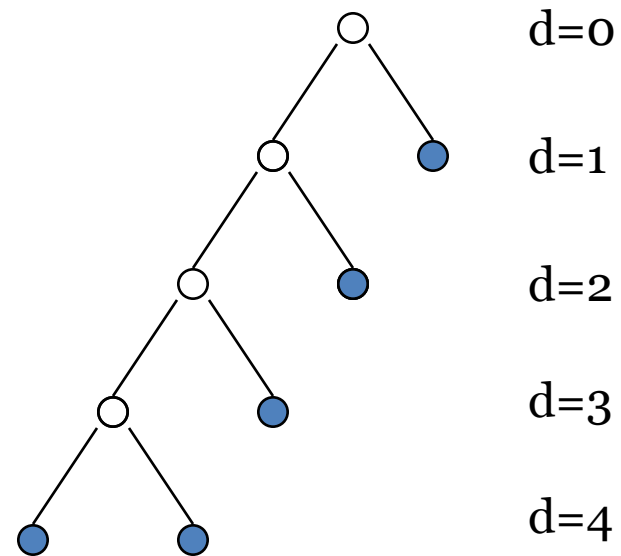


# What is the Complexity of Depth-First Search?

- Time Complexity
  - maximum tree depth =  $m$
  - assume (worst case) that there is 1 goal leaf at the RHS at depth  $d$  so DFS will generate  $O(b^m)$



- Space Complexity
  - how many nodes can be in the queue (worst-case)?
  - at depth  $m$  we have  $b$  nodes
  - and  $b-1$  nodes at earlier depths
  - total =  $b + (m-1)*(b-1) = O(bm)$



## Examples of Time and Memory Requirements for Depth-First Search

Assuming  $b=10$ ,  $m = 12$ , 10000 nodes/sec,  
1kbyte/node

Depth of Solution	Nodes Generated	Time	Memory
2	$10^{12}$	3 years	120kb
4	$10^{12}$	3 years	120kb
8	$10^{12}$	3 years	120kb
12	$10^{12}$	3 years	120kb

# Depth-First Search (DFS) Properties

- **Complete?**

- **No.** Not complete if tree has unbounded depth

- **Optimal?**

- **No.** Solution found first may not be the shortest possible

- **Time complexity?**

$$O(b^m)$$

- **Exponential** exponential in the maximum depth of the search tree  $m$

- **Space complexity?**

$$O(bm)$$

**linear** in the maximum depth of the search tree  $m$

- **Linear**

# Comparing DFS and BFS

- **Time complexity:** same, but
  - In the worst-case BFS is always better than DFS
  - Sometime, on the average DFS is better if:
    - many goals, no loops and no infinite paths
- BFS is much worse memory-wise
  - DFS is linear space
  - BFS may store the whole search space.
- In general
  - BFS is better if goal is not deep, if infinite paths, if many loops, if small search space
  - DFS is better if many goals, not many loops,
  - DFS is much better in terms of memory



Thank You!

**Any Questions?**